

# MAP Estimation Via Agreement on Trees: Message-Passing and Linear Programming

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## Outline



Introduction

- Maximum a Posteriori Probability and Graphical Model
- Definition and Symbols



#### Motivation

- Convexity of Phi
- Upper Bounds

Tree-Reweighted Message-Passing Algorithms 3 Max-Marginals Algorithm





Maximum a Posteriori Probability and Graphical Model Definition and Symbols

#### What Is This Paper About

- MAP derived from a graph is upper bounded by the linear combination of the sub-trees of the graph.
- Although the number of sub-trees may be untraceable, the problem turns out to be solvable using local marginal information that is only related to the nodes and edges in the dual space because of the convexity of the upper bound.
- There is a constraint imposed on the upper bound: the count of edges of the graph and the count of edges in the sub-trees need to be consistent.
- The constraint can be met with a special message passing construction: tree-reweighted belief propagation.



Maximum a Posteriori Probability and Graphical Model Definition and Symbols

## Markov Random Fields

- Advantage of MRF
  - Well structured isotropic behavior
  - Local dependencies
- Disadvantages of MRF
  - Difficult to compute probability
  - Parameter estimation is hard
- Applications: too many to list
  - Machine learning, Imaging, Computer Vision, Economics etc.



Maximum a Posteriori Probability and Graphical Model Definition and Symbols

#### Random Fields

- Random variables Y can be considered a Markov Random Field (MRF) on S if:
  - P(Y) > 0•  $P(y_i|y_{S-\{i\}}) = P(y_i|y_{N_i})$

We can also formulate this based on graphical model. Let G = (V, E) be a graph with vertices V and edges E. Vertices  $V = X \cup Y$  with X as the observation and Y (label) as random variables. The factorization is defined by: **Hammersley-Clifford's Theorem.** Assume that  $p(y_1, \ldots, y_n) > 0$  Then,

$$p(y) = \frac{1}{Z} exp\left(-\sum_{C} \phi_{C}(x_{C})\right)$$

where C is a clique, subset of nodes that are fully connected in the graph.



Maximum a Posteriori Probability and Graphical Model Definition and Symbols

#### Graphical Representation



Let's denote black nodes as observed nodes  $y_i$  and white nodes as hidden nodes  $x_i$ . The overall joint probability of p(x, y) is:

$$p(x,y) = \frac{1}{Z} \prod_{ij} \psi_{ij}(x_i, x_j) \prod_i \phi_i(x_i, y_i)$$

where the  $\psi_{ij}(x_i, x_j)$  and  $\phi_i(x_i, y_i)$  are joint compatibility functions. Then the Maximum A Posteriori (MAP) is given by:  $\operatorname{argmax}_x p(x|y)$ .



Maximum a Posteriori Probability and Graphical Model Definition and Symbols

# A Image Denoise/Segmentation Example



Restored Image (ICM)

Restored Image (Graph cuts)

Pattern Recognition and Machine Learning (C.M. Bishop)



Maximum a Posteriori Probability and Graphical Model Definition and Symbols

#### MAP and Graphical Model

- Message passing (MP) is popular in solving MAP problems (integer programming) in acyclic graph.
- MP exploits the conditional independent properties which is the key to factorize the graph.



$$p(x,y) = \frac{1}{Z} \prod_{i} \phi_i(x_i) \prod_{i,j} \psi_{ij}(x_i, x_j)$$



Maximum a Posteriori Probability and Graphical Model Definition and Symbols

#### MP in Undirected Graph

A typical message passing route is shown as:



The belief at a node *i* is proportional to the product of the local evidence at that node  $(\phi_i(x_i))$ , and all messages coming into node *i*:

$$b_i(x_i) = k\phi_i(x_i) \prod_{j \in \mathcal{N}(i)} m_{ji}(x_i)$$
  

$$m_{ij}(x_j) = \sum_{x_i} \phi_i(x_i)\psi_{ij}(x_i, x_j) \prod_{k \in \mathcal{N} \setminus j} m_{ki}(x_i)$$



Maximum a Posteriori Probability and Graphical Model Definition and Symbols

#### Preliminaries

An undirected graph is defined as G(V, E). For each  $s \in V$ , let  $X_s$  be a random variable taking values  $x_s$  in sample space  $\mathcal{X}_s$  and  $\mathcal{X}_s := \{0, \ldots, m_s - 1\}$ . For n = |V| elements, X takes values x in the Cartesian product space  $\mathcal{X}^n := \mathcal{X}_1 \times \mathcal{X}_2 \times \ldots \times \mathcal{X}_n$ . A full collection of potential functions associated with a given clique C is mapping  $\phi : \mathcal{X}^n \to \mathbb{R}^d$  with  $\{\phi_\alpha | \alpha \in \mathcal{I}\}$  and  $d = |\mathcal{I}|$ .  $\theta = \{\theta_\alpha | \alpha \in \mathcal{I}\}$  is the vector of parameter. Then, strictly positive MRF can be represented as:

$$\mathsf{p}(x; heta) \propto \exp\left\{ \langle heta, \phi(x) 
angle 
ight\} \equiv \exp \sum_{lpha \in \mathcal{I}} heta_lpha \phi_lpha(x)$$

For easy annotation, we define the following indicator functions:

$$\begin{aligned} \{\delta_j(x_s)|j \in \mathcal{X}_s\}, \text{ for } s \in V\\ \{\delta_j(x_s)\delta_k(x_t)|(j,k) \in (X)_s \times \mathcal{X}_t\}, \text{ for } (s,t) \in E \end{aligned}$$

Marginal distribution can be represented by:

$$\mu_{s;j} := \mathbb{E}_{p}[\delta_{j}(\mathsf{x}_{s})] = \sum_{x \in \mathcal{X}^{n}} p(x)\delta_{j}(\mathsf{x}_{s})$$
$$\mu_{st;jk} := \mathbb{E}_{p}[\delta_{j}(\mathsf{x}_{s})\delta_{k}(\mathsf{x}_{t})] = \sum_{x \in \mathcal{X}} p(x)[\delta_{j}(\mathsf{x}_{s})\delta_{k}(\mathsf{x}_{t})]$$



Maximum a Posteriori Probability and Graphical Model Definition and Symbols

#### Linear Constraints and MAP Estimation

Linear Constraints

$$\begin{split} &\sum_{j\in\mathcal{X}_{s}}\mu_{s;j}=1,\;\forall s\in V\\ &\sum_{(j,k)\in\mathcal{X}_{s}\times\mathcal{X}_{l}}\mu_{st;jk}=1,\;\forall (s,t)\in \textit{E},j\in\mathcal{X}_{s}\\ &\sum_{k\in\mathcal{X}_{l}}\mu_{st;jk}=\mu_{s;j},\forall (s,t)\in\textit{E},j\in\mathcal{X}_{s} \end{split}$$

#### MAP Estimation

Let  $\overline{\theta}$  be a given vector of parameter of  $\mathbb{R}^d$ . Let  $\overline{\theta}_s(x_s) := \sum_{j \in \mathcal{X}_s} \overline{\theta}_{s;j} \delta_j(x_s)$ . Let  $\overline{\theta}_{st}(x_s; x_t) := \sum_{(j,k) \in \mathcal{X}_s \times \mathcal{X}_t} \overline{\theta}_{st;jk} \delta_j(x_s) \delta_k(x_t)$ . MAP is to maximize:

$$\langle \overline{\theta}, \phi(x) \rangle := \sum_{s \in V} \overline{\theta}_s(x_s) + \sum_{(s,t) \in E} \overline{\theta}_{st}(x_s, x_t)$$

For the convenience, we define the MAP as follows:

$$\Phi_{\infty}(\overline{\theta}) := \max_{x \in (X)^n} \langle \overline{\theta}, \phi(x) \rangle$$
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Convexity of Phi Upper Bounds

## Convexity

Claim:  $\Phi_\infty$  is convex in terms of  $\overline{\theta}$ 

#### Proof

Let's start with a more general log-partition function:  $\Phi(\theta) = \log \sum_{x} \exp\{\theta^T \phi(x)\}$ .

$$\begin{aligned} \frac{\partial \Phi}{\partial \theta_k} &= \frac{\sum_x \exp\{\theta^T \phi(x)\}\phi_k(x)}{\sum_x \exp\{\theta^T \phi(x)\}} \\ &= \frac{\sum_x \exp\{\theta^T \phi(x)\}\phi_k(x)}{\exp \Phi(\theta)} \\ &= \sum_x \exp\{\theta^T \phi(x) - \Phi(\theta)\}\phi_k(x) \\ &= \sum_x p(x;\theta)\phi_k(x) \\ &= \mathbb{E}\{\phi_k(X)\} \\ \frac{\partial^2 \Phi}{\partial \theta_k \partial \theta_l} &= \frac{\partial}{\partial \theta_l} \sum_x \exp\{\theta^T \phi(x) - \Phi(\theta)\}\phi_kx \\ &= \sum_x \exp\{\theta^T \phi(x) - \Phi(\theta)\}[\phi_l(x) - \frac{\partial \Phi(\theta)}{\partial \theta_l}]\phi_k(x) \\ &= \mathbb{E}\{\phi_k(X)\phi_l(X)\} - \{\phi_k(X)\}\mathbb{E}\{\phi_l(X)\} \\ &= Cov_p\{\phi_k(x), \phi_l(x)\}. \end{aligned}$$



Convexity of Phi Upper Bounds

#### Illustration of Edge Appearance in Probabilities



For 
$$p(x,\overline{\theta}) \propto \exp(x_1x_2 + x_2x_3 + x_3x_4 + x_4x_1)$$
,  
 $\overline{\theta} = [0 \ 0 \ 0 \ 0 \ 1 \ 1 \ 1 \ 1]$ . Then  $\rho(T_i) = 1/4$  and  $\rho_e = 3/4$  for each  $e \in E$ . The can construct a member  $\theta$  as follows:

$$\begin{array}{rcl} \theta(T_1) &=& (4/3)[0\ 0\ 0\ 0\ 1\ 1\ 1\ 0] \\ \theta(T_1) &=& (4/3)[0\ 0\ 0\ 0\ 1\ 1\ 0\ 1] \\ \theta(T_1) &=& (4/3)[0\ 0\ 0\ 0\ 1\ 0\ 1\ 1] \\ \theta(T_1) &=& (4/3)[0\ 0\ 0\ 0\ 0\ 1\ 1\ 1] \end{array}$$

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Convexity of Ph Upper Bounds

#### Upper Bounds via Convex Combinations

Let  $\rho^i$  be a finite collection of nonnegative weights that sum to one, s.t.,  $\sum_i \rho^i \theta^i = \overline{\theta}$ . Applying Jensen's inequality yields the upper bound:  $\Phi_{\infty} \leq \sum_i \rho^i \Phi_{\infty}(\theta^i)$ . Each  $\Phi_{\infty}(\theta^i)$  represents an acyclic subgraph, for which exact computation are tractable. The index *i* corresponds to a spanning tree of the graph and all parameters are required to respect the structure of the tree. For a convex combination of tree parameters,  $\mathcal{E}_{\rho}[\theta(T)] := \sum_{T} \rho(T)\theta(T)$ :

#### Tightness of Upper Bounds

$$0 \leq \left[\sum_{T} \rho(T) \Phi_{\infty}(\theta(T))\right] - \Phi(\overline{\theta})$$
  
= 
$$\left[\sum_{T} \rho(T) \Phi_{\infty}(\theta(T))\right] - \langle \overline{\theta}, \phi(x^*) \rangle$$
  
= 
$$\sum_{T} \rho(T) \left[\Phi_{\infty}(\theta(T)) - \langle \theta(T), \phi(x^*) \rangle\right]$$

The bound is tight if and only if  $x^* \in \cap_T OPT(\theta(T))$  for some  $x^* \in OPT(\overline{\theta})$ .



Convexity of Phi Upper Bounds

## Objective

$$\begin{cases} \min_{\theta} & \sum_{T} \rho(T) \Phi_{\infty}(\theta(T)) \\ \text{s.t.} & \sum_{T} \rho(T) \theta(T) = \overline{\theta} \end{cases}$$

#### Theorem

The optimal value of the above problem can be obtained by:

$$\max_{\mu} \sum_{s \in V} \sum_{j} \mu_{s;j} \overline{\theta}_{s;j} + \sum_{(s,t) \in E} \sum_{j,k} \mu_{st;jk} \overline{\theta}_{st;jk}$$

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Convexity of Phi Upper Bounds

Claim: 
$$\Phi_{\infty}(\overline{\theta}) = \max \sum_{s \in V} \sum_{j} u_{s;j} \overline{\theta}_{s;j} + \sum_{(s,j) \in E} \sum_{j,k} \mu_{st;jk} \overline{\theta}_{st;jk}$$

#### Proof

By definition of  $\Phi_{\infty}$ ,  $\max_{x \in \mathcal{X}^n} \langle \overline{\theta}, \phi(x) \rangle = \max_{p \in \mathcal{P}} \sum_{x \in \mathcal{X}^n} p(x) \langle \overline{\theta}, \phi(x) \rangle$ .

$$\begin{split} \sum_{x \in \mathcal{X}^n} p(x) \langle \overline{\theta}, \phi(x) \rangle &= \sum_{x \in \mathcal{X}^n} p(x) \left\{ \sum_{s \in V} \overline{\theta}_s(x_s) + \sum_{(s,t) \in E} \overline{\theta}_{st}(x_s, x_t) \right\} \\ &= \sum_{s \in V} \sum_j \mu_{s:j} \overline{\theta}_{s:j} + \sum_{(s,t) \in E} \sum_{j,k} \mu_{st;jk} \overline{\theta}_{st;jk} \end{split}$$
  
where  $\mu_{s:j} := \sum_{x \in \mathcal{X}^n} p(x) \delta_j(x_s)$  and  $\mu_{st;jk} := \sum_{x \in \mathcal{X}^n} p(x) \delta_{jk}(x_s, x_t)$ 

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Convexity of Ph Upper Bounds

#### Lagrange Dual of the Objective Function

$$\mathcal{L} = \sum_{T} \rho(T) \Phi_{\infty}(\theta(T)) + \tau(\sum_{T} \rho(T)\theta(T) - \overline{\theta})$$
  
$$= \sum_{T} \rho(T) [\Phi_{\infty}(\theta(T)) - \langle \theta(T), \tau \rangle] + \langle \tau, \overline{\theta} \rangle$$

The Lagrange dual is then:

$$\sum_{T} \rho(T) \inf_{\theta(T)} [\Phi_{\infty}(\theta(T)) - \langle \theta(T), \tau \rangle] + \langle \tau, \overline{\theta} \rangle$$

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So,  $\frac{\partial \Phi_{\infty}(\theta(T)) - \langle \theta(T), \tau \rangle}{\partial \theta} = 0$  and  $\tau = E[\phi(x)]$ .



Max-Marginals Algorithm

### Max-Marginal Factorization

Any tree-structured distribution can be factorized in terms of its max-marginals as follows:

$$p(x;\theta(T)) \propto \prod_{s \in V} \mu_s(x_s) \prod_{(s,t) \in E(T)} \frac{\mu_{st}(x_s, x_t)}{\mu_s(x_s)\mu_t(t)}$$

The tree-structured parameter  $\theta(T)$  is defined in terms of logarithms of  $\mu$ :

$$\begin{array}{lcl} \theta_s^n(T)(x_s) & = & \log \mu_s(x_s) \; \forall s \in V \\ \theta_{st}^n(x_s, x_t) & = & \log \frac{\mu_{st}(x_s, x_t)}{\mu_s(x_s)\mu_t(x_t)} \; (s,t) \in E(T) \end{array}$$

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Max-Marginals Algorithm

#### Edge Based Reparameterization Updates

For iterations n = 0, 1, 2, ..., update the max-marginals as follows:

$$\mu_{s}^{n+1}(x_{s}) = k \mu_{s}^{n}(x_{s}) \prod_{t \in \Gamma(s)} \left[ \frac{\max_{x_{t}'} \mu_{st}^{n}(x_{s}, x_{t}')}{\mu_{s}^{n}(x_{s})} \right]^{\rho_{st}}$$
$$\mu_{st}^{n+1}(x_{s}, x_{t}) = k \frac{\mu_{st}^{n}(x_{s}, x_{t})}{\max_{x_{t}'} \mu_{st}^{n}(x_{s}, x_{t}') \max_{x_{s}'} \mu_{st}^{n}(x_{s}', x_{t})} \mu_{s}^{n+1}(x_{s}) \mu_{t}^{n+1}(x_{t})$$

In terms of messages, max-marginals are as follows:

$$\mu_{s}(x_{s}) \propto \phi_{s}(x_{s}) \prod_{v \in \Gamma(s)} [M_{vs}(x_{s})]^{\rho_{vs}}$$
  
$$\mu_{st}(x_{s}, x_{t}) \propto \phi_{st}(x_{s}, x_{t}) \frac{\prod_{v \in \Gamma(s) \setminus t} [M_{vs}(x_{s})]^{\rho_{vs}}}{[M_{ts}(x_{s})]^{(1-\rho_{ts})}} \times \frac{\prod_{v \in \Gamma(t) \setminus s} [M_{vt}(x_{t})]^{\rho_{vt}}}{[M_{st}(x_{t})]^{(1-\rho_{st})}}$$



Max-Marginals Algorithm

#### Edge Based Reparameterization Updates

#### The above construction establishes that for all $x \in \mathcal{X}^n$ , we have:

$$\sum_{T} \rho(T) \theta(T)(x) = \sum_{s \in V} \overline{\theta}_s(x_s) + \sum_{(s,t) \in E} \overline{\theta}_{st}(x_s, x_t)$$

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Max-Marginal Algorithm

#### Parallel Tree-Reweighted Message Passing Algorithm



Initialize the message m<sup>0</sup> = m<sup>0</sup><sub>ij</sub> with arbitrary positive numbers.
 for each iteration, update the message as follows:

$$m_{ts}^{n+1}(x_s) = k \sum_{x'_t} \exp\left(\frac{1}{\rho_{ts}} \mu_{st}(x_s, x'_t) + \mu_t(x'_t)\right) \times \frac{\prod_{v \in \Gamma(t) \setminus s} [m_{vt}^n(x'_t)]^{\rho_{vt}}}{[m_{st}^n(x'_t)]^{(1-\rho_{st})}}$$



Max-Marginals Algorithm

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#### Local Beliefs

Once the process has converged, the local belief can be calculated as:

$$b_s(x_s) = k \exp(-\mu_s(x_s)) \prod_{t \in \Gamma(s)} [m_{ts}(x_s)]^{
ho_{ts}}$$



Max-Marginals Algorithm

#### Testing Example





Max-Marginals Algorithm

## Testing Example





## Conclusion

- The paper provides an upper bound for the optimal (MAP) configuration.
- The new tree-reweighted free energy is convex with respect to the max-marginals vector
- No sufficient conditions to guarantee convergence on graphs with cycles.
- The demo code can be downloaded from http://code.google.com/p/random-field-blief-propagation/